Introduction OpenMP Parallelization

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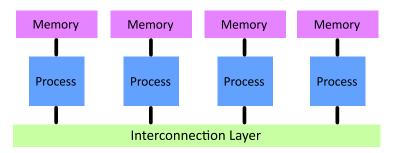
June 2, 2011

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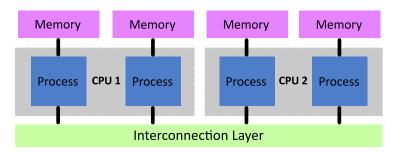
 Shared-Memory vs. Distributed-Memory Parallelism: Processes and Threads

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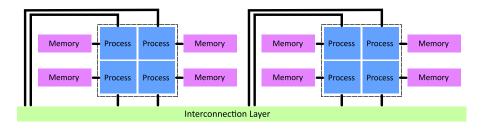
- Approaches to Shared-Memory Parallelization
- OpenMP: Overview
- OpenMP: A Practical Introduction
- Hybrid Parallelism



MPI-type distributed-memory parallelism.

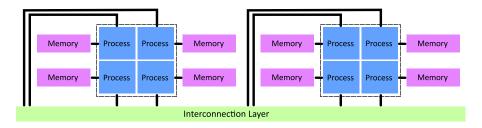


- MPI-type distributed-memory parallelism
- ► Multicore CPU → multiple MPI processes per CPU. Interconnection via networking software stack.
- ► Multiple CPUs per node → Interconnection via networking software stack.
- ► Between nodes: Interconnect via networking hardware.



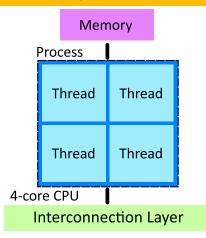
- ► Iowa's Helium cluster has 2 quad-core CPUs per node.
- MPI: each core runs one process with its own memory. Communication via network stack within node and with other nodes.

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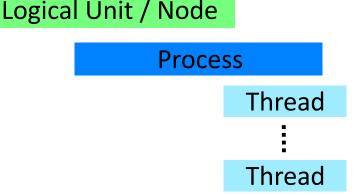
- ► Iowa's Helium cluster has 2 quad-core CPUs per node.
- MPI: each core runs one process with its own memory. Communication via network stack within node and with other nodes.
- This seems like a great waste! Why not share the memory within a CPU or even a node and bypass the interconnect?

Threaded Shared-Memory Parallelism



- A process may have multiple parallel threads sharing the same memory (multi-threading). Each process has at least one thread.
- One thread per physical core (note: Intel Hyper-Threading 2 virtual cores per physical core)

Parallelism Hierarchy



- Only one thread communicates with the "outside".
- ► At least one thread per process (exactly 1 in classical MPI setup).
- ► No shared memory between typical cluster nodes.
- Number of cores per node keeps increasing.

Architectures for Shared-Memory Parallelism





Multi-core laptop

SGI Altix UV "Blacklight" at PSC

- Any shared memory symmetric multiprocessor machine (SMP).
 → Any modern laptop/desktop; any one cluster compute node. Limited to physical unit (cluster node).
- ► Non-Uniform Memory Access machines system-wide shared memory architectures → 1000s of cores.

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 - Code unchanged.
 - Compiler specific. Results vary greatly. Won't parallelize complex loops.
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 - Requires major code re-write from single-thread version. Lots of pedestrian work.

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OpenMP

- ► Full control, run-time thread allocation.
- Only small code changes needed.
- Convenient, high-level interface.

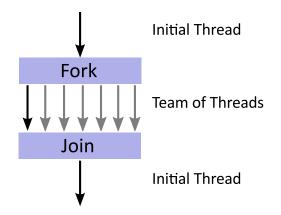
Introduction: What is OpenMP?



http://www.openmp.org

- OpenMP us an application programming interface (API) for shared-memory multi-threading.
- OpenMP is not an external library. It is implemented directly by the compiler.
- ► OpenMP 1.0 in 1997; Current: OpenMP 3.0 (3.1 coming).
- OpenMP works with C/C++ and Fortran.

Basic: The OpenMP Concept



 Fork-Join multi-threading model for mixing serial with shared-memory-parallel program sections.

Basics: Compiling with OpenMP

 The compiler must be told to use OpenMP. This is accomplished via compiler flags that differ between compilers.

```
GNU gcc/gfortran (open source)
gcc -fopenmp ..., g++ -fopenmp ...,
gfortran -fopenmp ...
```

Intel Compilers

```
icc -openmp ..., icpc -openmp ...,
ifort -openmp ...
```

Basics: OpenMP API Components

OpenMP has three basic components:

Pre-processor directives.

C/C++: #pragma omp parallel ... F90 : !\$OMP PARALLEL ... !\$OMP END PARALLEL

Runtime library routines.

C/C++: #include <omp.h>

F90: use module omp_lib

 Environment variables. To be set in the shell, e.g., <u>OMP_NUM_THREADS=8</u>
 A first OpenMP Program in Fortran 90:

program omp1 use omp_lib implicit none

```
!$OMP PARALLEL ! Fork threads
write(6,*) "my thread id: ", omp_get_thread_num()
!$OMP END PARALLEL ! Join threads
```

end program omp1

A First OpenMP Program in C

```
A first OpenMP Program in C:
#include <stdio.h>
#include <omp.h>
int main(void) {
 #pragma omp parallel // Fork threads
  ſ
  printf("my thread id: %d\n", omp_get_thread_num());
  }
                         // Join threads
}
```

In this example (and in the previous Fortran one), every thread redundantly executes the code in the parallel region.

Useful OpenMP Library Functions

- omp_get_thread_num: current thread index (0, 1, . . .)
- omp_get_num_threads: size of the active team
- omp_set_num_threads: set size of the thread team (make this call outside of a parallel region)
- omp_get_max_threads: maximum number of threads
- omp_get_num_procs: number of cores available

There are a couple more – see the OpenMP reference manual for a full list and description.

Parallelizing Loops

Finally doing something useful...

The compute-intense parts of most codes are loops over large datasets that carry out many floating point operations.

```
do k=1,nz
  do j=1,ny
   do i=1,nx
    [do something crazy complicated]
    enddo
   enddo
enddo
enddo
```

[do something crazy complicated] is executed nx*ny*nz times!

Basic Worksharing: Parallel do/for

Parallelizing "do/for" loops:

► C for loop

```
#pragma omp parallel
#pragma omp for
for(i=0;i<n;i++) {
   // do something in parallel
}</pre>
```

or, using a combined directive:

```
#pragma omp parallel for
for(i=0;i<n;i++) {
   // do something in parallel
}
```

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Basic Worksharing: Parallel do/for

Parallelizing "do/for" loops:

```
> Fortran do loop
    !$OMP PARALLEL
    !$OMP D0
    do i=1,n
        ! do something in parallel
    enddo
    !$OMP END D0
    !$OMP END PARALLEL
```

```
or, using a combined directive:
```

```
!$OMP PARALLEL DO
do i=1,n
  ! do something in parallel
enddo
!$OMP END PARALLEL DO
```

Basic Worksharing with Loops: Rules

- Only standard for/do loops can be parallelized.
 while loops cannot.
- Program correctness must not depend upon which thread executes a particular iteration. For example:

Does not work:	Works:
x(1) = 0	do i=1,n
do i=2,n	x(i) = (i-1)*f
x(i) = x(i-1) + f	enddo
enddo	

 Branching statements such as break, exit, continue, goto, return etc. are not allowed.

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A somewhat mindless example:

Let's write a simple code that fills an array of length n with numbers and see how this can be sped up with OpenMP.

```
program omp2ser
  implicit none
  integer :: i
  integer, parameter :: n = 26000000
  real*8,allocatable :: myarray(:)
  allocate(myarray(n))
  do i=1,n
   myarray(i) = 5*i**3 + i**2 + i + sin(1.0*i)**2
  enddo
  deallocate(myarray)
```

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end program omp2ser

Basic Worksharing: Basic loop example - parallel

```
program omp2
  implicit none
  integer :: i
  integer, parameter :: n = 26000000
  real*8,allocatable :: myarray(:)
  allocate(myarray(n))
  !$OMP PARALLEL DO
  do i=1,n
    myarray(i) = 5*i**3 + i**2 + i + sin(1.0*i)**2
  enddo
  !$OMP END PARALLEL DO
  deallocate(myarray)
end program omp2
```

Let's try this on helium:

```
export OMP_NUM_THREADS=1;time ./omp2
export OMP_NUM_THREADS=2;time ./omp2
export OMP_NUM_THREADS=4;time ./omp2
export OMP_NUM_THREADS=8;time ./omp2
```

 \rightarrow Won't see major improvement at more than 8 threads, since helium only has 8 physical cores.

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Let's reduce the problem size n by a factor of 100 and try again: export OMP_NUM_THREADS=1;time ./omp2b100 export OMP_NUM_THREADS=2;time ./omp2b100 export OMP_NUM_THREADS=4;time ./omp2b100

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 \rightarrow No improvement with increasing number of threads! Why?

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Let's reduce the problem size n by a factor of 100 and try again:

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export OMP_NUM_THREADS=4;time ./omp2b100

export OMP_NUM_THREADS=8;time ./omp2b100

 \rightarrow No improvement with increasing number of threads! Why? OpenMP "fork/join" process requires time. If problem size too small, forking/joining dominates compute time.

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In more complex settings, it becomes necessary to tell OpenMP what variables are **private** to each thread and which are **shared**.

By default, all variables are assumed to be **shared**. Exceptions: Loop counters of the outermost loop and variables declared inside the parallel region (only in C).

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Declaring vars private/shared:

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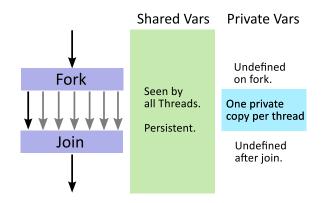
By default, all variables are assumed to be **shared**. Exceptions: Loop counters of the outermost loop and variables declared inside the parallel region (only in C).

Declaring vars private/shared:

```
→ In C:
#pragma omp parallel for private(pvar1,pvar2) shared(svar)
→ In Fortran:
!$OMP PARALLEL DO PRIVATE(pvar1,pvar2) SHARED(svar)
[...]
```

!\$OMP END PARALLEL DO

Basic Worksharing: Private/Shared Vars



- Shared vars: Seen by all threads, but not more than one thread must write to a shared var at a time. Persistent.
- Private vars: Private "copy" for each thread. Undefined when the thread team is created; undefined after parallel region.

Consider this code snippet:

```
[...]
!$OMP PARALLEL DO
do i=1,n
    x = 5*i**3 + i**2 + i + sin(1.0*i)**2
    myarray(i) = x
enddo
!$OMP END PARALLEL DO
[...]
```

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i is private, but n, x, and myarray are shared.

Consider this code snippet:

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enddo
!$OMP END PARALLEL DO
[...]
```

i is private, but n, x, and myarray are shared. data race condition: x is updated inconsistently and uncontrollably by multiple threads!

Private/Shared Example: fixed

```
Fixed:
[...]
!$OMP PARALLEL DO PRIVATE(x)
do i=1,n
    x = 5*i**3 + i**2 + i + sin(1.0*i)**2
    myarray(i) = x
enddo
!$OMP END PARALLEL DO
[...]
```

i and x are private. n, myarray are shared. Outside the parallel segment, i and x are undefined.

Another Loop Example

Suppose we wanted to parallelize

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```
[...]

sum = 0.0d0

do i=1,n

val = f(i)

sum = sum + val

enddo

[...]
```

Another Loop Example

Suppose we wanted to parallelize

```
[...]

sum = 0.0d0

do i=1,n

val = f(i)

sum = sum + val

enddo

[...]
```

First attempt:

```
[...]
sum = 0.0d0
!$OMP PARALLEL DO PRIVATE(val)
do i=1,n
  val = f(i)
  sum = sum + val
enddo
!$OMP END PARALLEL DO
[...]
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Another Loop Example

Suppose we wanted to parallelize

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[...]

sum = 0.0d0

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enddo

[...]
```

First attempt:

```
[...]
sum = 0.0d0
!$OMP PARALLEL DO PRIVATE(val)
do i=1,n
val = f(i)
sum = sum + val
enddo
!$OMP END PARALLEL DO
[...]
```

Problem: Race condition in the updating of sum!

Another Loop Example: fixed (1: CRITICAL)

One way of fixing this is the **!\$OMP** CRITICAL directive:

```
[...]
sum = 0.0d0
!$OMP PARALLEL DO PRIVATE(val)
do i=1,n
  val = f(i)
  !$OMP CRITICAL
  sum = sum + val
  !$OMP END CRITICAL
enddo
!$OMP END PARALLEL DO
[...]
```

The CRITICAL directive ensures that only one thread accesses sum at a time.

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An even better way of dealing with this issue is a **sum** reduction:

```
[...]
sum = 0.0d0
!$OMP PARALLEL DO PRIVATE(val) REDUCTION(+:sum)
do i=1,n
  val = f(i)
  sum = sum + val
enddo
!$OMP END PARALLEL DO
[...]
```

The **REDUCTION** clause tells OpenMP that the team of threads must safely add to sum so that it assumes the same value as in the serial case.

Numerical reduction operators: +, -, * See reference manual/sheet for more.

Advanced Stuff: Ask Google!

Some more useful clauses that modify OpenMP behavior

- NO WAIT clause don't wait after a loop (or other directive) inside a parallel section until all threads are done.
- SCHEDULE(STATIC) evenly divide iterations of a loop among threads.
- SCHEDULE(DYNAMIC[, chunk]) divide work into chunk-sized parcels. If a thread is done with a chunk, it grabs another one. Default chunk size is 1.
- SCHEDULE(GUIDED[, chunk]) divide work into chunks of exponentially decreasing size. chunk is the minimum chunk size. Default is 1.

Advanced Stuff: Ask Google!

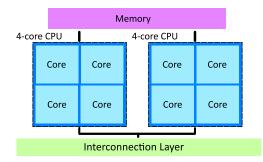
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There are a more OpenMP directives.

- ► SECTIONS non-iterative work sharing.
- ► BARRIER force threads to wait for each other.
- ► ORDERED force sequential order in a loop.
- ► MASTER section in a loop executed only by the master.
- SINGLE section in a loop executed only by one thread.

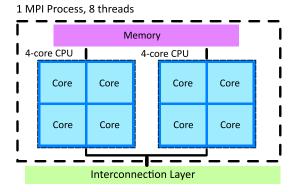
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 Modern cluster supercomputers have nodes with an increasing number of cores. Helium: 8 cores per node (two 4-core CPUs).

All cores within a node share the same main memory.

Hybrid Parallelism



Hybrid Parallelism:

- Node-local OpenMP.
- Internode MPI.
- Reduces communication overhead. Optimal number of MPI processes per node depends on software & hardware.